

# Development of a Translator from LLVM to ACL2\*

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In our current work a library of formally verified software components is to be created, and assembled, using the Low-Level Virtual Machine (LLVM) intermediate form, into subsystems whose top-level assurance relies on the assurance of the individual components. We have thus undertaken a project to build a translator from LLVM to the applicative subset of Common Lisp accepted by the ACL2 theorem prover. Our translator produces executable ACL2 formal models, allowing us to both prove theorems about the translated models as well as validate those models by testing. The resulting models can be translated and certified without user intervention, even for code with loops, thanks to the use of the `def : : ung` macro which allows us to defer the question of termination. Initial measurements of concrete execution for translated LLVM functions indicate that performance is nearly 2.4 million LLVM instructions per second on a typical laptop computer. In this paper we overview the translation process and illustrate the translator’s capabilities by way of a concrete example, including both a functional correctness theorem as well as a validation test for that example.

## 1 Introduction

In our current work, we need to create formally verified software systems from a library of verified components assembled using the Low-Level Virtual Machine (LLVM) intermediate form [13]. To accomplish this we have undertaken a project to build a translator from LLVM to the applicative subset of Common Lisp [12] accepted by the ACL2 theorem prover [10], and perform verification of the software system using ACL2’s automated reasoning capabilities.

LLVM is the intermediate form for many common compilers, including the `clang` compiler used by Mac OS X and iOS developers. LLVM supports a number of language frontends, and LLVM code generation targets exist for a wide variety of machines, including both CPUs and GPUs. LLVM is a register-based intermediate language in Static Single Assignment (SSA) form [4]. As such, LLVM supports any number of registers, each of which is only assigned once, statically (dynamically, of course, a given register can be assigned any number of times). Andrew Appel has observed that “SSA form is

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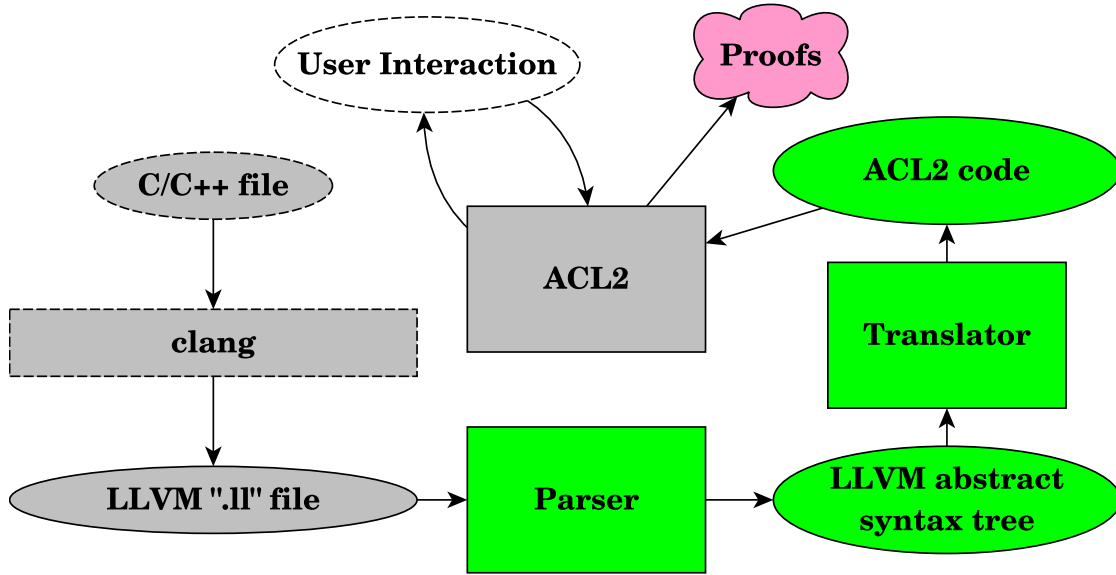


Figure 1: LLVM-to-ACL2 translation toolchain.

a kind of functional programming” [1]; this observation, in turn, inspired us to build a translator from LLVM to the applicative subset of Common Lisp accepted by the ACL2 theorem prover. Our translator produces an executable ACL2 specification that is able to support proof-based verification, as well as validation via testing.

## 2 Toolchain Overview

Our translation toolchain architecture is shown in Figure 1. The left side of the figure depicts a typical compiler frontend producing LLVM intermediate code. LLVM output can be produced either as a binary “bitcode” (.bc) file, or as text (.ll file). We chose to parse the text form, producing an abstract syntax tree (AST) representation of the LLVM program. Our translator converts the AST to an ACL2 model of the code which ACL2 is able to certify automatically. Once certified, the model can be loaded together with conjectures that one wishes to prove about the code. In addition to proving theorems about the translated LLVM code, ACL2 can also be used to execute test vectors against the translated model at reasonable speeds.

The translator is written in OCaml [5]. OCaml was chosen because the translator developer was proficient in that language. The translator comprises less than 9500 lines of OCaml code, and employs a parser generator whose input file, describing the LLVM grammar, is some 1800 lines. The translator successfully parses all 5000+ legal .ll files in the LLVM source distribution. The translator produces an AST from the input, removes aliases, extracts functions from labelled basic blocks, constructs parameter lists, determines declaration order, then generates the ACL2 code for each function.

```

unsigned long occurrences(unsigned long val, unsigned int n,
                          unsigned long *array) {
    unsigned long num_occur = 0;
    unsigned int j = 0;
    for (j = 0; j < n; j++) {
        if (array[j] == val) num_occur++;
    }
    return num_occur;
}

```

---

Figure 2: Example C code to count occurrences of an input value in an array.

### 3 An Example

As an example, consider the C source code of Figure 2. This function counts the number of occurrences of a given value in the first  $n$  elements of an array. (NB: By default the `clang` compiler treats all `int` values as 32 bits wide, and all `long` values as 64 bits wide.)

This is one of the more simple examples in the regression test suite for our framework, and it fails to exercise many of the more advanced features of our translator and proof infrastructure. However, its relative simplicity allows us to narrate a complete translation and verification within the confines of this paper.

We produce the LLVM code for this function by invoking `clang` as follows: `clang -O4 -S -emit-llvm occurrences.c`. The generated LLVM code for `clang` version 4.2 (which supports LLVM 3.2) is excerpted in Figure 3.

Observe that LLVM output is similar to assembly code, with labels and low-level opcodes like `br` (branch), `icmp` (integer compare) and `load` (load from memory). Registers are prepended with the “%” character, and are given sometimes-meaningful names. Consistent with the SSA philosophy, no register appears on the left hand side of an assignment (“=”) more than once. A peculiar feature of LLVM code is the `phi` instruction, which provides register renaming at a branch target. We will use the `phi` in our ACL2 translation to match formal to actual parameters, as will be detailed later.

### 4 Translation to ACL2

Treating the SSA as a functional program, we convert register assignments into nested `let` bindings. Each label is treated as a unique function, so we produce `defun` forms for `@occurrences`, `.lr.ph`, and `._crit_edge`. The formal parameters for the functions can be determined by consulting the left hand side of the `phi` functions; thus, `._crit_edge` should have a formal parameter `%num_occur_dot_0_dot_lcssa` in its parameter list (accounting for the differences in allowed characters in parameter names). We also need to identify parameters that are read, but not modified — most of our functions will thus require `%val`, `%n` and `%array` as input parameters.

We are left, then, with the question of how to translate memory and memory transactions. Typically in ACL2, a machine state data structure is declared, and passed as a parameter to all functions that read and/or write elements of the state. If a given function updates the state, the modified state must be returned. Obviously, for a large state, functional update of the state can become quite expensive. In an

```

define i64 @occurrences(i64 %val, i32 %n, i64* %array) {
  %1 = icmp eq i32 %n, 0
  br i1 %1, label %._crit_edge, label %.lr.ph

.lr.ph:
  %indvars.iv = phi i64 [ %indvars.iv.next, %.lr.ph ], [ 0, %0 ]
  %num_occur.01 = phi i64 [ %num_occur.0, %.lr.ph ], [ 0, %0 ]
  %2 = getelementptr inbounds i64* %array, i64 %indvars.iv
  %3 = load i64* %2, align 8, !tbaa !1
  %4 = icmp eq i64 %3, %val
  %5 = zext i1 %4 to i64
  %num_occur.0 = add i64 %5, %num_occur.01
  %indvars.iv.next = add nuw nsw i64 %indvars.iv, 1
  %lftr.wideiv = trunc i64 %indvars.iv.next to i32
  %exitcond = icmp eq i32 %lftr.wideiv, %n
  br i1 %exitcond, label %._crit_edge, label %.lr.ph

._crit_edge:
  %num_occur.0.lcssa = phi i64 [ 0, %0 ], [ %num_occur.0, %.lr.ph ]
  ret i64 %num_occur.0.lcssa
}

```

---

Figure 3: LLVM code for the occurrences example.

earlier version of the translator [9], we utilized an ACL2 single-threaded object (stobj) [2] to represent state. The destructive update property of stobjs provided good performance when executing translated functions on concrete state. However, stobjs are not currently compatible with the deferred termination technology that we wished to employ, so stobjs were abandoned in favor of typed records [6].

#### 4.1 LLVM State Representation

The raw LLVM state is represented using the following `defstructure` (`defstructure` can be found in `data-structures/structures.lisp` in the ACL2 community books):

```
(defstructure raw-st
  (retval (:assert (natp retval) :rewrite ))
  (stack  (:assert (usb32 stack) :rewrite ))
  (frame  (:assert (usb32 frame) :rewrite ))
  (mem    (:assert (wf-mem mem) :rewrite ))
  (:options :guards
            (:predicate stp)
            (:keyword-updater update-raw-st)))
```

This declaration indicates that the LLVM state is composed of four fields: `retval`, `frame`, `stack`, and `mem`.

The LLVM memory model assumes a 32-bit underlying architecture. In other words, the primitive `int` type is 32 bits wide, as are memory addresses. The `mem` field contains a memory of 8-bit words, implemented as an association list. While the size of the memory array is, technically, unbounded, as a practical matter, its size is limited to 4 GB since we use 32-bit addresses. The memory could easily be upgraded to 64-bit addressing in the future. Little-endian byte ordering is assumed by the low-level `rd.n` and `wr.n` primitives, but big-endian could also be readily supported.

The `retval` field is used to return values from procedure calls. This was needed historically since all translated functions accept state and return only state. In the future this field could be replaced by a stack lookup or a multiple-valued return. Note that the size of `retval` is unbounded (a `natp`). This allows us to return objects such as structures by converting them to natural numbers.

The `stack` field is a 32-bit pointer that identifies the next available stack location in memory. The stack grows towards infinity. The `frame` field is a 32-bit pointer that identifies the top of the previous stack frame.

#### 4.2 Completing the Translation

The translated function for the LLVM code beginning at label `._crit_edge` is depicted below. The translator ensures that the state record, `st`, is passed to, and returned from, all translated functions, even simple ones such as this. Referring to the LLVM of Figure 3, `._crit_edge` does nothing more than stash the value of `num_occur` in the `retval` field of the state and return the updated state.

```
(def::un occurrences_%_dot__crit_edge (%num_occur_dot_0_dot_lcssa st)
  (declare (xargs :signature ((i64_p stp) stp)))
  (let* ((st (update-retval %num_occur_dot_0_dot_lcssa st)))
    st))
```

This definition employs features of Greve’s `def` package, provided as part of the ACL2 community books [7] [8]. The `def::un` macro, found in the `coi/util/defun` book, improves upon ACL2 `defun` by providing both input and output “type” signatures. In the example above, the signature

```
(declare (xargs :signature ((i64_p stp) stp)))
```

says that the given function takes two inputs, the first of which satisfies the `i64_p` predicate and the second of which satisfies the `stp` predicate, and returns one output, satisfying the `stp` predicate. The `:signature` form accepts all standard Common Lisp type declarations (i.e.: `string`, `integer`, `unsigned-byte`) as well as generic predicate symbols and even lambda expressions. In the `defun` form generated by `def::un` the input element types become guards and the signature declaration is interpreted as a theorem about the return type of the function, in this case:

```
(DEFTHM I64_P-STP-IMPLIES-STP-OCCURRENCES_%_DOT__CRIT_EDGE
  (IMPLIES (AND (I64_P %NUM_OCCUR_DOT_O_DOT_LCSSA)
                (STP ST))
    (STP (OCCURRENCES_%_DOT__CRIT_EDGE
          %NUM_OCCUR_DOT_O_DOT_LCSSA ST))))
:RULE-CLASSES
(:REWRITE
 (:FORWARD-CHAINING
  :TRIGGER-TERMS ((OCCURRENCES_%_DOT__CRIT_EDGE
                    %NUM_OCCUR_DOT_O_DOT_LCSSA ST)))))
```

This theorem states that, for a given invocation of the function, if the inputs satisfy their associated predicates, then the value returned by the function satisfies its predicate. This rule is written as a forward-chaining rule so that this “type” information is readily propagated. Such properties are often needed and `def::un` saves us the drudgery of explicitly writing rules of this form ourselves for each new function that we define.

It should be noted that `def::un` is similar in spirit to the `defunc` (“defun with contracts”) form in ACL2s [14], although the details of the respective contract specifications differ (`defunc` features `:input-contract` and `:output-contract` keywords, whose interpretation is somewhat different from the `:signature` of `def::un`). The `def::un` macro also allows a user to specify congruence relations satisfied in each argument position, a capability which `defunc` does not provide.

### 4.3 Translating Loops

Loops in LLVM are translated into recursive functions in ACL2. Automatically generating recursive functions in ACL2 is challenging because ACL2 requires the identification of a well-founded relation for use as a measure to ensure that the recursion terminates. Without user input, however, automatically identifying such a measure for arbitrary LLVM code is impossible. To alleviate this issue we employ the `def::ung` macro. This macro, found in `coi/defung/defung.lisp` in the ACL2 community books, allows us to admit arbitrary recursive functions without the need of a measure. Rather, the macro generates a companion domain predicate\* which, when true, ensures that the function terminates. We thus exchange the need to identify a measure before admitting the function for the inconvenience of reasoning about the function’s domain predicate during subsequent proofs. As an aside, the `def::ung` macro

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\*The macro also generates a measure function which is guaranteed to decrease with each recursive call in the domain of the function.

supports the same set of `:signature` declarations as `def::un` and also produces function definitions that can be executed efficiently. The `def::ung` macro, however, does not currently support stobjs or multiple-value returns, limitations that have driven several trade-offs in our translation framework.

The original C source in our example (see Figure 2) used a `for` loop, which can be naturally expressed as a `while` loop (test at the top of the loop), but the corresponding loop in LLVM (the `.lr.ph` block) is of `do-while` form (test at the bottom of the loop). Transforming loops into `do-while` form is standard for `clang` in all but the least optimized mode of operation. The high-level specifications used in our work, however, employ `while` loops exclusively. To ease the process of relating our LLVM implementations to their high-level specifications, our LLVM translator has been designed to emit only `while` loops.

Another useful optimization is to isolate the potentially complex body of a loop from the recursive call in the ACL2 output. Hence for any given loop in LLVM, we emit a single “step” function containing the body of the loop. This “step” function is called by the `while` function in ACL2 to advance the state. During proof, the step function can oftentimes be disabled, thus simplifying and speeding the proof process.

Upon completion of a loop the next action performed is typically a jump to a subsequent block of sequential code. Having the base case of a recursive function invoke an arbitrarily complex block of code, however, complicates the reasoning process unnecessarily. In our translation, therefore, when the recursive function terminates, it simply returns the live variables and the current state. We introduce a “`while_wrap`” function that calls the `while` loop and then calls the function corresponding to the next LLVM block to be executed.

The general form of the generated ACL2 code for the “Nth” loop of an arbitrary LLVM procedure named “`fun`” is shown in Figure 4. Note that the `mvlist` macro returns multiple values as a single list and the `metlist` macro emulates multiple-value binding for functions returning such lists of values. These macros can be found in the book `coi/util/mv-nth`. We use these macros to skirt the fact that `def::ung` does not support multiple-value returns.

Observe that, although one could never accuse this clique of generated functions of being terribly efficient (e.g., there are many more function calls and returns than one would like to see, with much marshalling/unmarshalling of data), nonetheless this structure preserves the tail recursive nature of the LLVM input, allowing us to perform computations on large input terms without exhausting the Lisp stack. We will revisit performance issues in Section 5.

The generated “step” function for the occurrences example is depicted in Figure 5. Note that the last thing that this step function does is return a list of parameters to be used for the next execution of the step function. Note particularly that the first value in the returned list, corresponding to the position of the `done` parameter, is set to the value of the `%exitcond` register, which indicates whether the updated loop index variable is equal to the array size. Also note that throughout the remainder of the example, we use the `bits` function from `rtl/re19` in the ACL2 community books whenever we need a result modulo some finite number of bits.

The top-level occurrences function, shown in Figure 6, is defined in terms of `occurrences_0`, the final function admitted by our loop translation. `occurrences_0` represents the start of the `do-while` loop.

## 5 Concrete Execution

It is advantageous to be able to validate the translated models by running them against concrete inputs. Since all of our functions are executable, we can readily perform such validation testing. In the ACL2

```

(def::un fun_continue (... st)
  (declare (xargs :signature ((... stp) stp)))
  (let ((st <post-loop functionality>))
    st))

(def::un fun_step_N (... st)
  (let ((st <loop functionality>))
    (let ((done <set or clear done bit>))
      (mvlist done ... st))))

(def::ung fun_step_N_while (done ... st)
  (declare (xargs :signature ((natp ... stp) ... stp)))
  (if (= done 1) (mvlist ... st)
      (metlist ((done ... st) (fun_step_N ... st))
                (fun_step_N_while done ... st))))

(def::un fun_step_N_while_wrap (... st)
  (declare (xargs :signature ((... stp) stp)))
  (metlist ((... st) (fun_step_N_while 0 ... st))
            (let ((st (fun_continue ... st)))
              st)))

(def::un fun_N (... st)
  (declare (xargs :signature ((... stp) stp)))
  (let ((done <set or clear done bit>))
    (if (= done 1) (fun_continue ... st)
        (fun_step_N_while_wrap ... st))))

```

---

Figure 4: Outline of the generated ACL2 code for an LLVM loop.



```

(def::un occurrences_step_0 (done %num_occur_dot_01 %indvars_dot_iv
                             %array %n %val st)
  (declare (xargs :signature ((natp i64_p i64_p _30_p i32_p i64_p stp)
                              natp i64_p i64_p _30_p i32_p i64_p stp)))
  (let*
    ((%2 (+ %array (_30_gep (list %indvars_dot_iv))))
     (%3 (i64_frombytes (loadbytes *i64_size* %2 st)))
     (%4 (icmp= %3 %val))
     (%5 (zext %4 1 64))
     (%_dot_num_occur_dot_0
      (bits (+ %5 %num_occur_dot_01) 63 0))
     (%indvars_dot_iv_dot_next
      (bits (+ %indvars_dot_iv (bits 1 63 0)) 63 0))
     (%lftr_dot_wideiv (bits %indvars_dot_iv_dot_next 31 0))
     (%exitcond (icmp= %lftr_dot_wideiv %n)))
    (mvlist %exitcond %_dot_num_occur_dot_0 %indvars_dot_iv_dot_next
            %array %n %val st)))

```

---

Figure 5: Generated ACL2 code for the occurrences loop body.

```

(def::un occurrences (%val %n %array st)
  (declare (xargs :signature (( i64_p i32_p _24_p stp) stp)))
  (let*
    ((st (init-stack-frame st))
     (st (begin-stack-frame st))
     (st (occurrences_0 %val %n %array st)))
    (end-stack-frame st)))

```

---

Figure 6: Generated top-level driver code for the translated occurrences example.

```

(def::un occurrences-test1 ()
  (declare (xargs :signature (() natp)))
  (let*
    ((myst (st 0 #xffff0000 #xffff0000 nil))
     (myst (update-mem
              (wr.n 8 #x8038 399
                (wr.n 8 #x8030 234
                  (wr.n 8 #x8028 0
                    (wr.n 8 #x8020 75
                      (wr.n 8 #x8018 399
                        (wr.n 8 #x8010 399
                          (wr.n 8 #x8008 (1- (expt 2 64))
                            (wr.n 8 #x8000 20 (mem myst)))))))))) myst)))
    (retval (occurrences 399 8 #x8000 myst))))

```

---

Figure 7: Concrete test case for the translated occurrences example.

code of Figure 7, we set up an initial state, in which the stack and frame pointers are set to a high memory location so that the stack cannot overwrite our data array (we place the array at address #x8000). We then write a number of 64-bit values into memory at increasing addresses, initializing the array. This state is then passed to the translated top-level occurrences function, along with inputs for the `val` and `n` parameters, as well as the location of the array in memory. After the execution of occurrences, the return value is fetched using the `retval` accessor.

As we have written the value 399 into the array three times, when we run `(occurrences-test1)` from the ACL2 prompt, it returns the correct value: 3.

We need not, however, restrict ourselves to small arrays. The underlying typed record representation for memory does not explicitly store “0” values; thus, any address that has not been explicitly written with a non-zero value is assumed to have a value of 0. Thus, we can replace the last line in the test of Figure 7 by

```
(retval (occurrences 0 1000000 #x8000 myst))
```

This corresponds to one million executions of the occurrences LLVM inner loop, plus some start-up and clean-up code (which we ignore here). Execution of this modified test returned the correct result (999993), and ACL2’s `time$` function returned 3.8 seconds of real time for the above on a late 2012 MacBook Pro. Further, the occurrences LLVM inner loop consists of 9 LLVM instructions; thus,

$$(1,000,000 * 9) / 3.8 \approx 2,370,000 \text{ LLVM instructions per second.}$$

This result is encouraging, considering that the translated Lisp code has not yet been engineered for performance in any serious way.

## 6 Reasoning about Translated Functions

In order to reason about a function such as `occurrences` in ACL2, we first need to perform abstraction on the data types; particularly, we wish to abstract the input array to a Lisp list. This can be done with the aid of a “lift” function, as follows:

```

(def::ung liftlist (done j array n st)
  (declare (xargs :signature ((natp natp natp natp stp) nat-listp)))
  (if (equal done 1) nil
      (let*
        ((ptr (+ array (* j 8)))
         (val (wfrombytes 8 (loadbytes 8 ptr st)))
         (j (bits (1+ j) 63 0))
         (done (if (equal (bits j 31 0) n) 1 0)))
        (cons val (liftlist done j array n st)))))

```

Note that we are modeling our lift function in a manner that closely reflects the behavior of our underlying LLVM implementation. In particular, we have chosen to admit our lift function using `def::ung` rather than to admit it as a standard ACL2 definition with a measure. While a reasonable measure for this function does, in fact, exist, the measure for a similar function that, for example, traversed a linked list rather than simply iterating through an array, would likely not exist. Thus, we proceed with this example as though a measure for our function is not available.

The list-based specification of occurrences is fairly conventional:

```

(def::un occurlist (val list)
  (declare (xargs :signature ((natp nat-listp) natp)))
  (if (endp list)
      0
      (+ (if (= val (car list)) 1 0)
         (occurlist val (cdr list)))))

```

We wish to prove that the translated occurrences function operating over an array in memory produces a result equal to the occurlist function operating over a (lifted) list. In order to carry out this proof, we must reason about the domain predicates generated by `def::ung`, which can be quite tedious. For example, we need for the domain of `occurrences_step_0_while` to be the same as the domain of `liftlist`. Unfortunately, however, `occurrences_step_while` has an additional argument. In order for the domains to be the same, we need to show that this argument is irrelevant. We can show this by induction, but first we need an appropriate induction scheme. This scheme is illustrated in `occurrences_step_0_while-induction-2`. Note that the admission of this scheme leverages the `occurrences_step_0_while-domain` and `occurrences_step_0_while-measure` functions introduced by `def::ung`.

```

(defun occurrences_step_0_while-induction-2 (done num num2 %next
                                             %array %n %val st)
  (declare
    (xargs :measure
      (occurrences_step_0_while-measure done num %next
                                           %array %n %val st)))
  (if (not (occurrences_step_0_while-domain done num %next
                                             %array %n %val st))
      st
      (if (equal done 1)
          (mvlist num %next %array %n %val st)

```

```

(metlist ((done2 num2 %next2 %array2 %n2 st2)
          (occurrences_step_0 done num2 %next
                               %array %n %val st))
  (metlist ((done num %next %array %n %val st)
            (occurrences_step_0 done num %next
                               %array %n %val st))
    (occurrences_step_0_while-induction-2 done num num2 %next
                                           %array %n %val st))))))

```

With the induction scheme in hand the proof of the irrelevance of the second argument is trivial.

```

(defthm domains-equiv-num-occur-irrelevant-lemma
  (implies
    (occurrences_step_0_while-domain done num_occur1
                                      j array n val st)
    (occurrences_step_0_while-domain done num_occur2
                                      j array n val st))
  :hints
  (("Goal" :induct
    (occurrences_step_0_while-induction-2 done num_occur1 num_occur2
                                           j array n val st)
    :in-theory (enable occurrences_step_0))))

```

Knowing that the second argument is irrelevant allows us to show that `occurrences_step_0_while-domain` and `liftlist-domain` are equivalent by mutual implication.

```

(defthm occurrences_step_0_while-implies-liftlist
  (implies (occurrences_step_0_while-domain done num_occur
                                              j array n val st)
    (liftlist-domain done j array n st))
  :hints (("Goal" :induct (occurrences_step_0_while done num_occur
                                              j array n val st)
    :in-theory (enable occurrences_step_0))))

(defthm liftlist-implies-occurrences_step_0_while
  (implies (not (occurrences_step_0_while-domain done num_occur
                                              j array n val st))
    (not (liftlist-domain done j array n st)))
  :hints (("Goal" :induct (liftlist done j array n st)
    :in-theory (enable occurrences_step_0))))

```

Once we know that the domains of the two functions are equivalent we can prove that the LLVM loop satisfies its specification. ACL2 is able to prove this theorem automatically because of the way

that we constructed the `liflist` function. (NB: the macro `val` in this example extracts the number of occurrences computed by the LLVM loop from the returned list of values.)

```
(defthm occurrences_rec_equiv--thm
  (implies
    (and (stp st) (natp done) (natp n) (natp array)
          (natp val) (bvecp val 64)
          (natp j) (natp num_occur) (bvecp num_occur 64))
    (equal (val 0 (occurrences_step_0_while done num_occur
                                              j array n val st))
            (bits (+ num_occur
                     (occurlist val (liflist done j array n st)))
                  63 0)))
  :hints (("Goal" :induct
            (occurrences_step_0_while done num_occur
                                       j array n val st)
            :in-theory
            (enable occurrences_step_0_while occurrences_step_0))))
```

The top-level specification is expressed in terms of `occurlist`:

```
(def::un occurrences_spec (val n array st)
  (declare (xargs :signature ((natp natp natp stp) natp)))
  (if (zp n)
      0
      (bits (occurlist val (liflist 0 0 array n st)) 63 0)))
```

The final equivalence theorem is then as follows:

```
(defthm occurrences_equiv--thm
  (implies (and (stp st) (natp n) (natp array) (bvecp val 64))
    (equal (retval (occurrences val n array st))
            (occurrences_spec val n array st))))
```

This result follows easily once we have proved `occurrences_rec_equiv--thm`.

## 7 Related Work

Zhao *et al.* [16] produced several different formalizations of operational semantics for LLVM in Coq [3], noting that their intention is to produce a verified LLVM compiler, similar to the verified CompCert compiler due to Leroy [11] (CompCert does not utilize the LLVM intermediate form). As such, their emphasis on formalizing LLVM operational semantics makes sense. We also considered creating an

“LLVM interpreter” in ACL2 (Zhao *et al.* utilized the OCaml extraction capability of the Coq environment to produce such an interpreter), resulting in a “deep embedding”, but decided that a translation to ACL2 (thus producing a “shallow embedding”) would allow us to begin proving properties about LLVM programs with much less effort. Our approach was also influenced by Magnus Myreen’s “decompilation into logic” work [15]. Our approach could be characterized as a sort of decompilation into logic, but we do not go to the same lengths as Myreen to assure that the decompilation process is sound. We also have the advantage of starting with a form that is functional, whereas Myreen has tackled the much more difficult problem of decompiling imperative machine code.

## 8 Conclusion and Future Work

We have built a translator from the LLVM intermediate form to the applicative subset of Common Lisp accepted by the ACL2 theorem prover. The translator produces executable tail-recursive ACL2 specifications, and we have utilized this executability in order to validate our translated models via testing. Performance of concrete execution for translated LLVM functions has been measured at nearly 2.4 million LLVM instructions per second on a typical laptop computer.

Future work will focus on continued enhancements to the translator and support books for enhanced automated reasoning, particularly dealing with translations that produce mutual recursions. The translator will also need to be updated to support new LLVM releases.

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